

Securify: Practical Security Analysis of Smart Contracts

<https://securify.ch>

Dr. Petar Tsankov

Scientific Researcher, ICE center, ETH Zurich

Co-founder and Chief Scientist, ChainSecurity AG

<http://www.ptsankov.com/>

[@ptsankov](#)

Inter-disciplinary and **inter-department** research center at ETH Zurich



Prof. Martin
Vechev



Prof. Laurent
Vanbever



Dr. Petar
Tsankov



Dr. Dana
Drachsler



Timon
Gehr



Ahmed
El-Hassany



Maria
Apostolaki



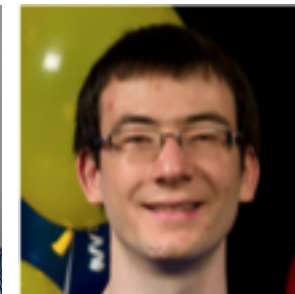
Rüdiger
Birkner



Samuel
Steffan



Roland
Meier



Johannes
Kapfhammer

Research @ ICE



Programmable networks



Blockchain security



Safe and interpretable AI



Security and privacy

Research @ ICE



Programmable networks



Blockchain security



Safe and interpretable AI



Security and privacy

What is a Smart Contract?

```
mapping(address => uint) balances;  
  
function withdraw() {  
    uint amount = balances[msg.sender];  
    msg.sender.call.value(amount);  
    balances[msg.sender] = 0;  
}
```

Transfer ETH to
the caller

- Small programs that handle cryptocurrencies
- Written in high-level languages (e.g., Solidity, Vyper)
- Executed on the blockchain (e.g. Ethereum)
- Usually no patching after release

What can happen when programs handle billions worth of USD?

Smart Contract **Security Bugs** in the News



Wallet bug freezes more than \$150 million worth of Ethereum



Breaking News: Bug Discovered In ICON (ICX) Smart Contract – Token Transfers Disabled

Critical Bug Found in ICON Smart Contract

By Oualou Awan-Momanyi Last updated Jun 16, 2018

NEWS ACTION NEWS



2 days ago

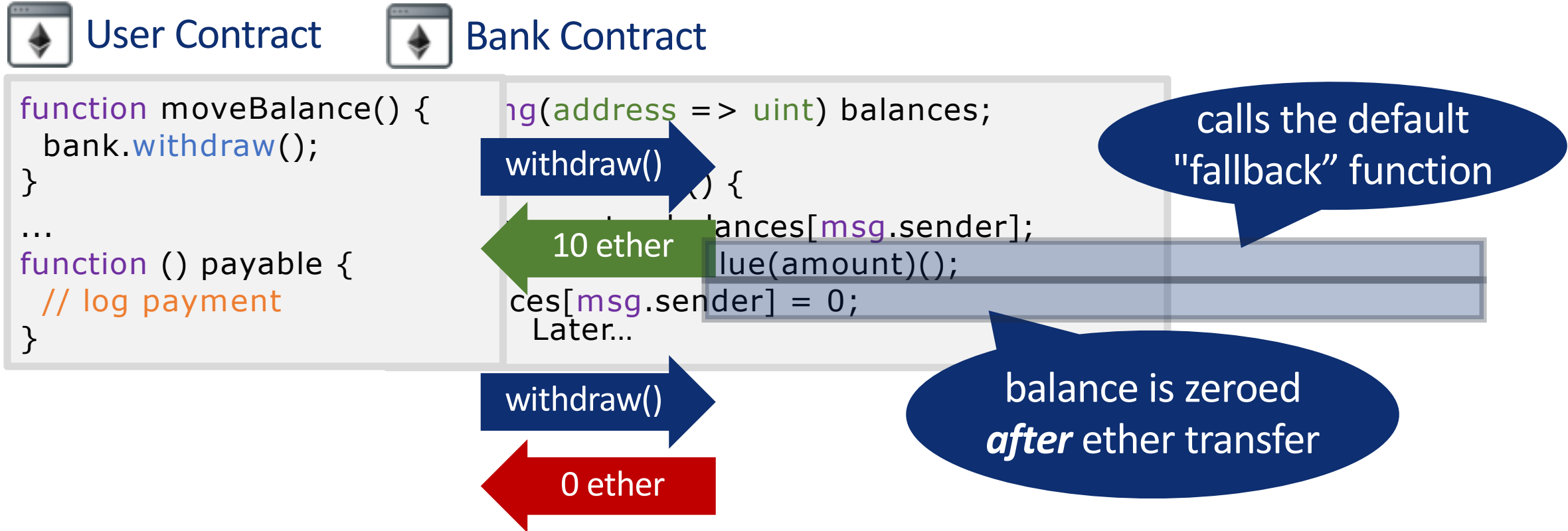


million worth of ethereum another hacker attack



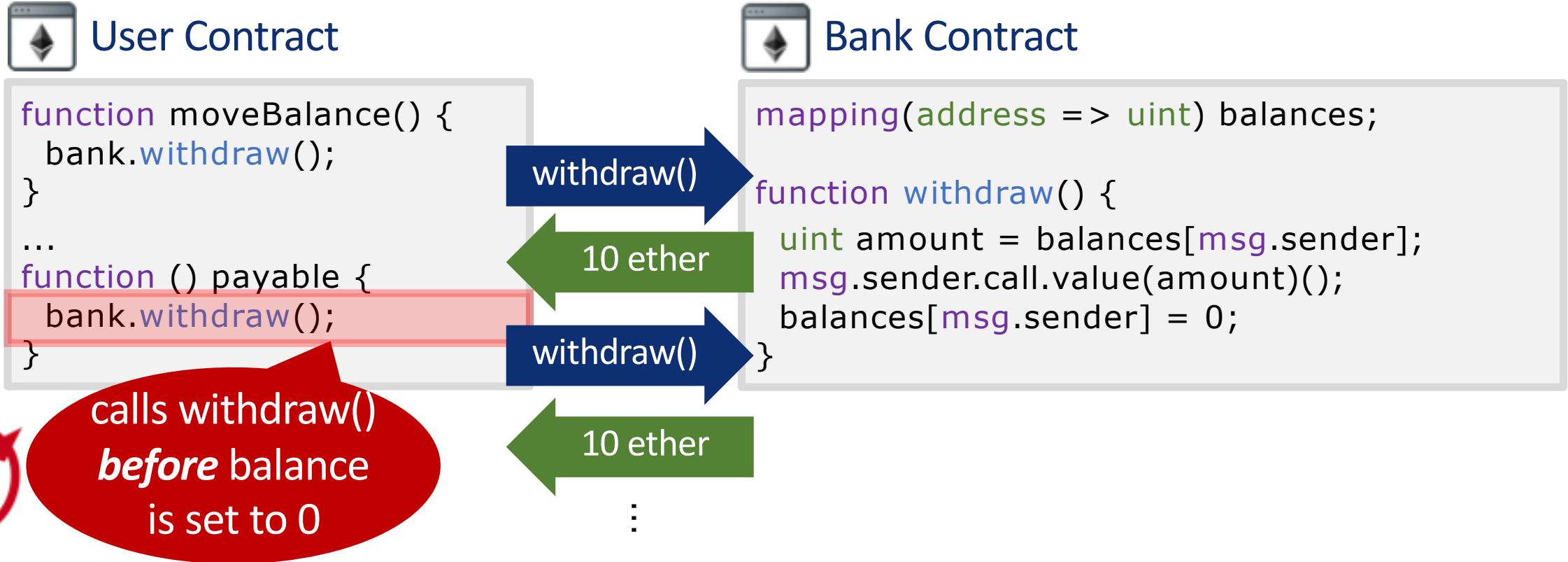
June 2016: The DAO hack

The DAO hack : Reentrancy



Can the user contract withdraw more than its balance?

The DAO hack: Reentrancy



An attacker used this bug to steal 3.6M ether (> **1B USD today**)

July 2017: Parity Multisig Bug 1

Parity Multisig **Bug 1**: Unprivileged Write to Storage



Any user may
change the
wallet's owner



Wallet Contract

```
address owner = ...;

function initWallet(address _owner) {
    owner = _owner;
}

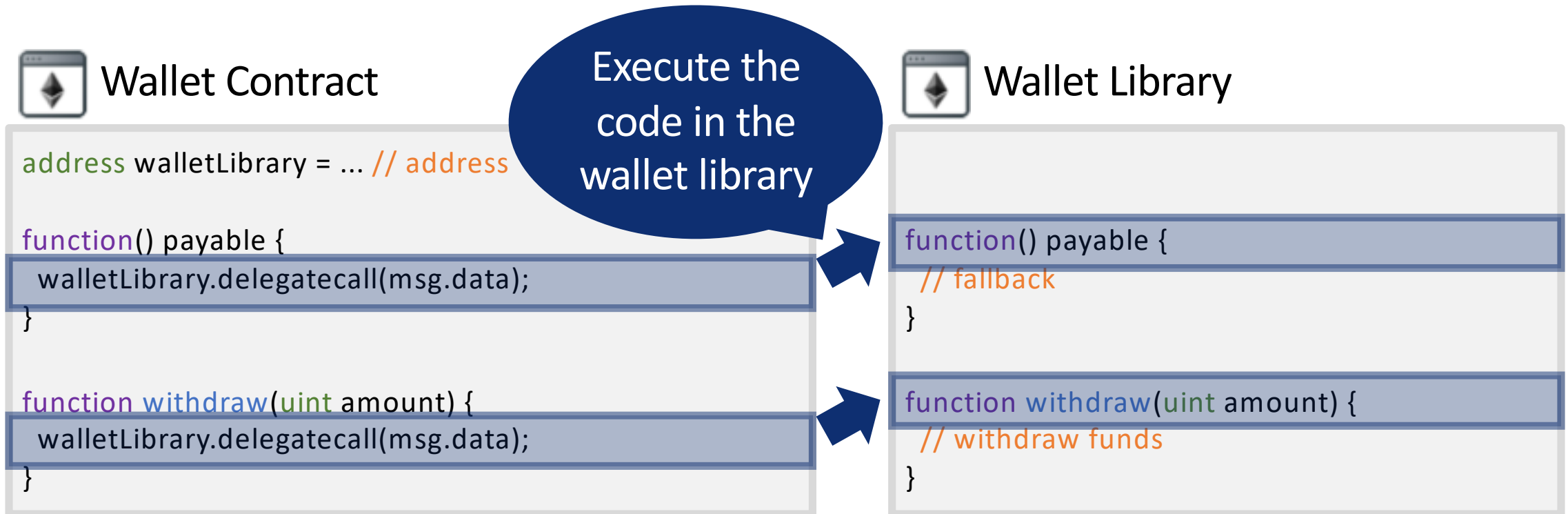
function withdraw(uint _amount) {
    if (msg.sender == owner) {
        msg.sender.transfer(_amount);
    }
}
```

Only the owner can
withdraw ether

An attacker used a similar bug to **steal \$30M** in July 2017

Four months later... Parity Multisig Bug 2

Parity Multisig **Bug 2**: Frozen Wallets



However, in Ethereum, smart contracts can be killed!

Parity Multisig **Bug 2**: Frozen Wallets



An attacker deleted the library



Wallet Contract

```
address walletLibrary = ... // address

function() payable {
  walletLibrary.delegatecall(msg.data);
}

function withdraw(uint amount) {
  walletLibrary.delegatecall(msg.data);
}
```

No withdraws are possible



Wallet Library

```
...

function() payable {
  // fallback
}

function withdraw(uint amount) {
  // withdraw funds
}
```



A user **froze \$170M** by deleting the wallet library

Relevant Security Properties...



Unexpected ether flows



Insecure coding, such as unprivileged writes



Use of unsafe inputs (e.g., reflection, hashing, ...)



Reentrant method calls (e.g., DAO bug)



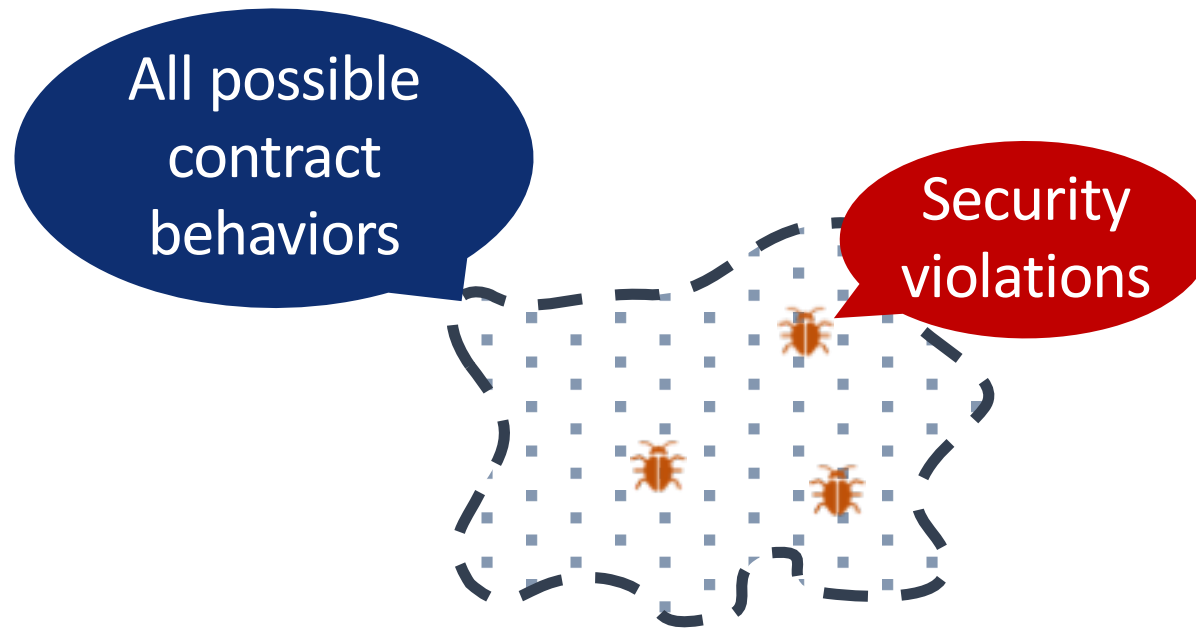
Manipulating ether flows via transaction reordering

Many of these are nontrivial trace-/hyper-properties



Automated Security Analysis of Smart Contracts: Challenges and Gaps

Security Analysis (high-level view)

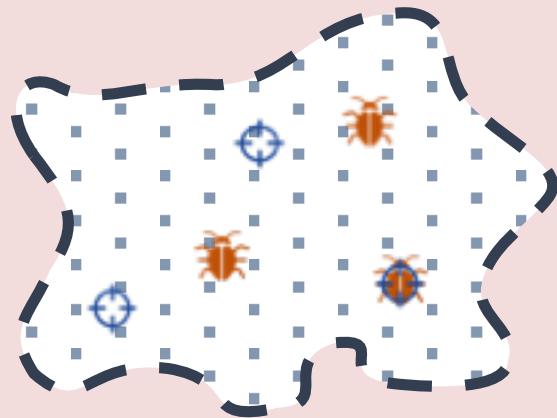


Minor issue 😊 : Smart contracts are written in Turing-complete languages

Automated Security Solutions



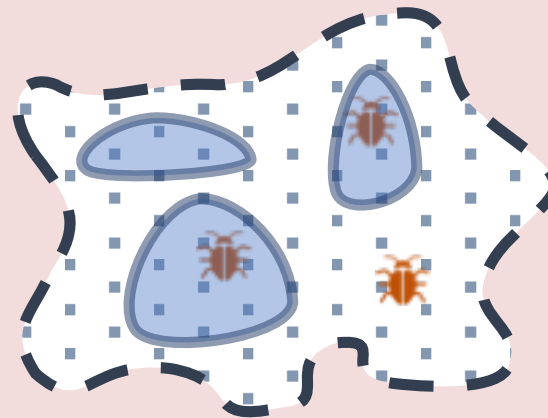
Truffle



Testing

Report true bugs
Can miss bugs

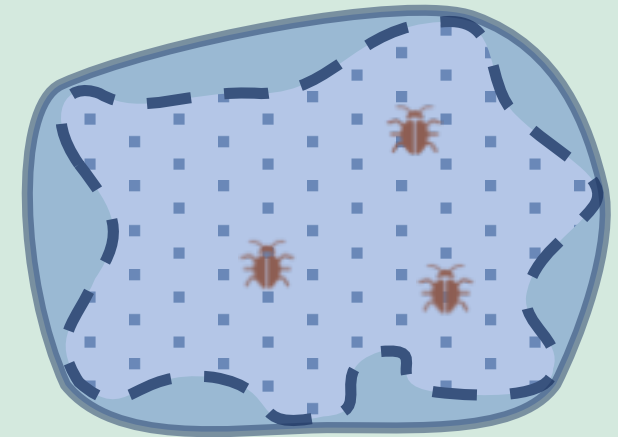
Oyente, Mythril, MAIAN



Dynamic (symbolic) analysis

Report true bugs
Can miss bugs

WANTED: Automated Verifier



Can report false alarms
No missed bugs

Bug finding

Verification

Domain-Specific Insight:

When contracts satisfy/violate a property, they often also satisfy/violate a much simpler property

Example: The DAO Hack

Security property

No state changes after call instructions

Hard to verify
in general

```
function withdraw() {  
  uint amount = balances[msg.sender];  
  msg.sender.call.value(amount)();  
  balances[msg.sender] = 0;  
}
```

Compliance pattern

No writes to storage **may follow** call instructions

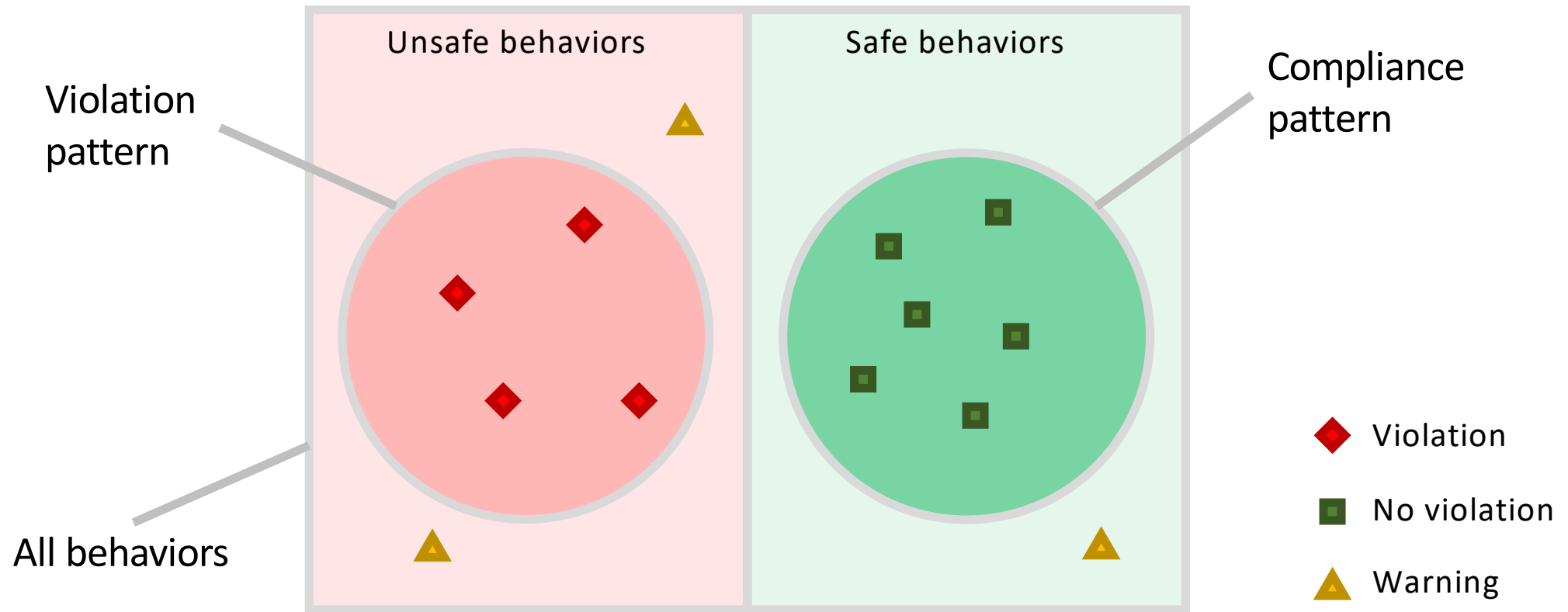
Verifies 91% of all
deployed contracts

Violation pattern

A write to storage **must follow** call instructions

Easier to check
automatically

Classifying Behaviors using Compliance and Violation Patterns



All unsafe behaviors are reported



A practical *verifier* for Ethereum smart contracts:

- fully-automated
- extensible
- scalable
- precise
- publicly available

www.securify.ch




The image features a dark blue background with a complex geometric pattern of thin, light blue lines and small dots. A horizontal band of a slightly lighter blue color runs across the middle of the image. The word "DEMO" is centered within this band in a white, sans-serif font.

DEMO

Beta version released in Fall 2017


- Regularly used by auditors to perform professional security audits

New release coming up very soon


-  95% positive feedback
-  > 8K uploaded smart contracts
-  > 800 users signed up for updates



Interesting discussions on Reddit

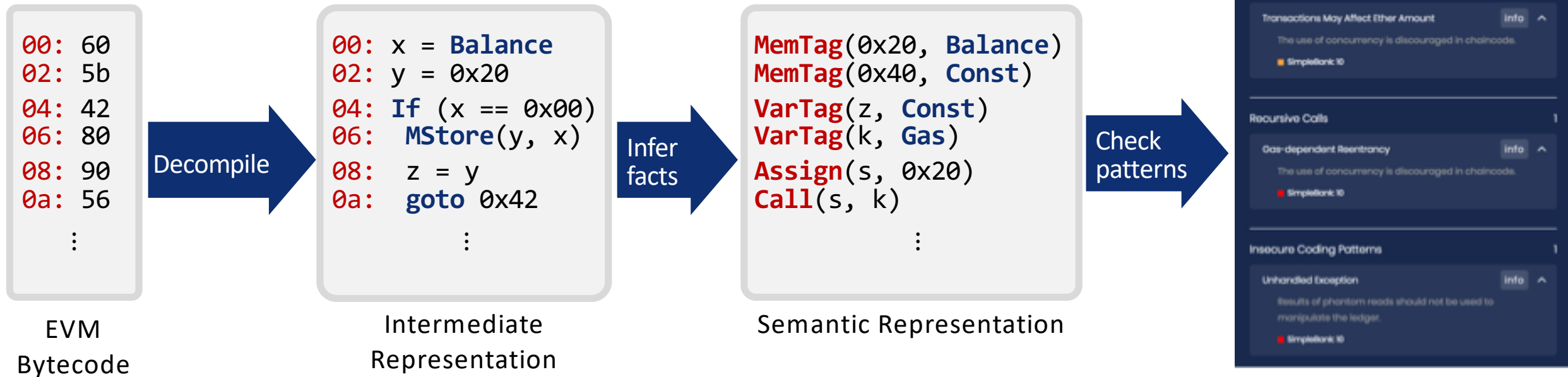
 [-] mcgravier 22 points 12 days ago
Seems almost too good to be true :) What are the limitations and how exactly does it work under the hood?

It's great that the authors of the tool are aware they are not capturing a full set of behaviors in the growing direction. That's the way to ensure safety properties without false-negatives. I'm interested how they compare their EVM semantics against other EVM implementations in the wild.

 [-] AlexanderSupersloth 12 points 12 days ago
Please, someone, humour a layman: how can a Turing complete language be formally verified?

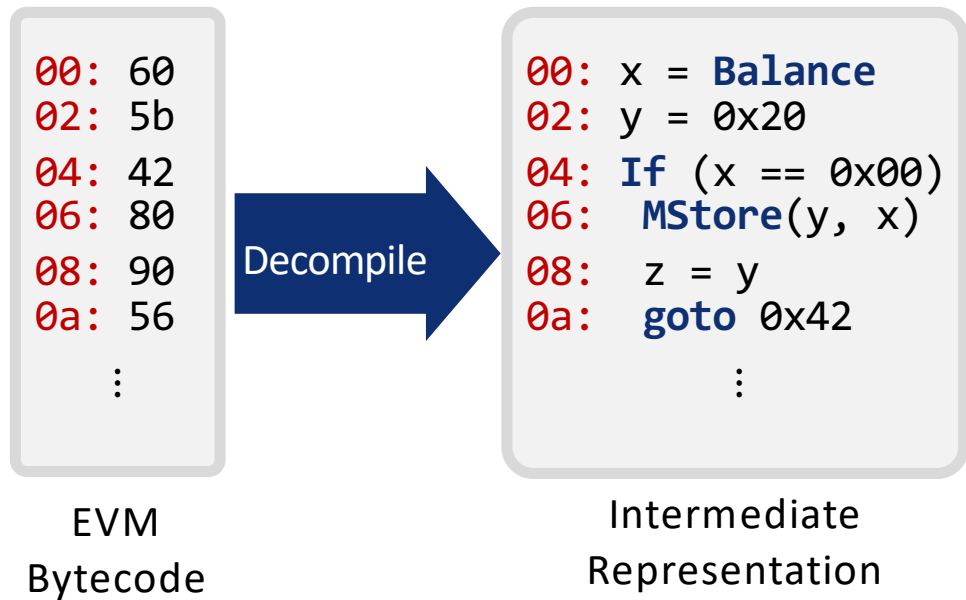
I thought formally verifiable languages were necessarily not Turing complete, and we can therefore not formally verify Solidity.

Securify: Under the Hood

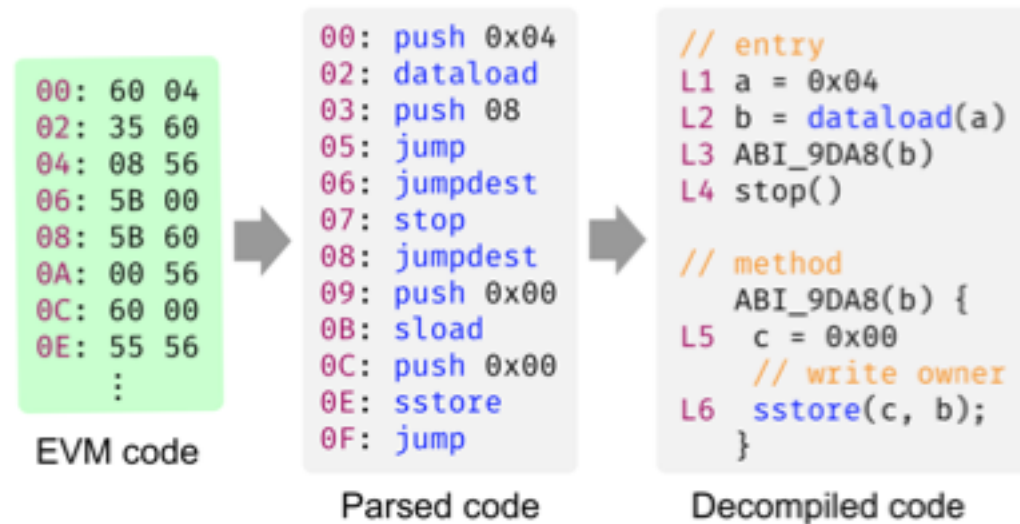


Fully automated, sound, scalable, extensible

Securify: Under the Hood



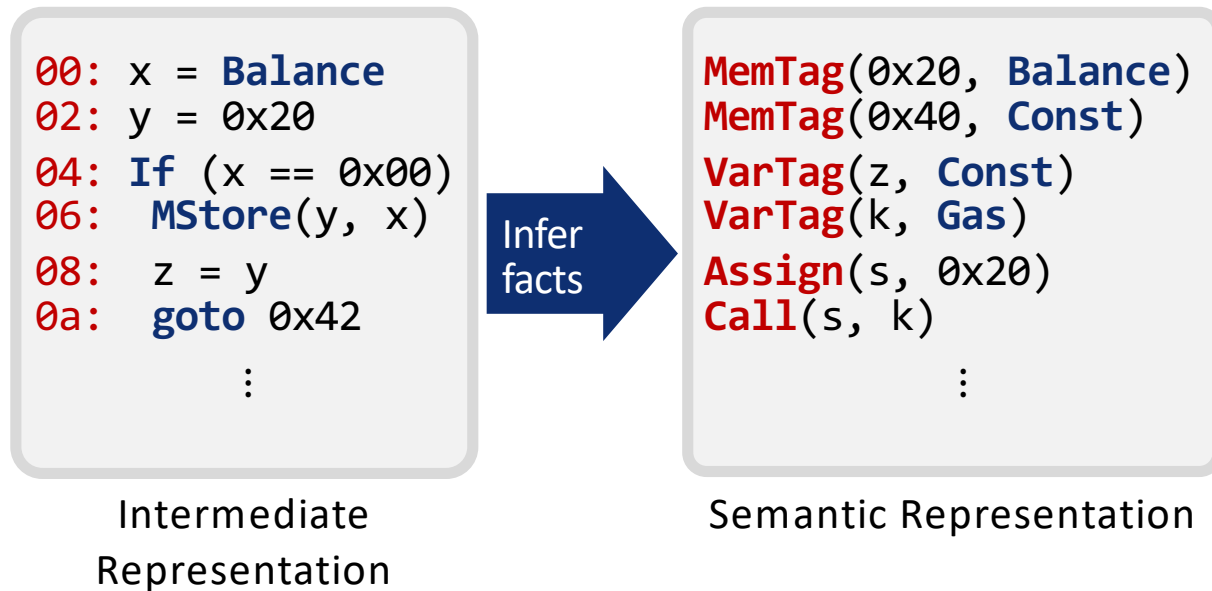
From EVM to CFG over SSA



Decompiling EVM bytecode:

- Convert into **static single assignment form** (each variable is assigned once)
- Perform **partial evaluation** (to resolve jump destination, memory/storage offsets)
- Identify and inline methods (to enable context-sensitive analysis)
- Construct **control-flow graph**

Securify: Under the Hood



Which facts are relevant for verifying smart contracts?

Semantic Facts

Many properties can be checked on the contract's dependency graph

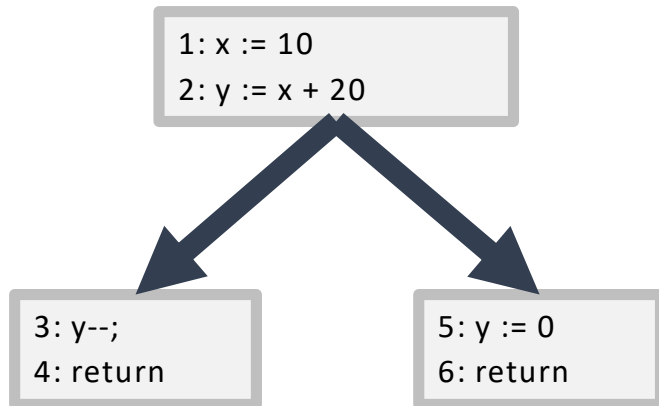
Flow dependencies	
$MayFollow(l, l')$	The instruction at label l may follow that at label l'
$MustFollow(l, l')$	The instruction at label l must follow that at label l'
Data dependencies	
$MayDepOn(x, t)$	The value of x may depend on tag t
$DetBy(x, t)$	For different values of t the value of x is different.

A tag can be an instruction (e.g. Caller) or a variable

The inference of all semantic facts is declaratively specified in Datalog

Example: *MayFollow*

$MayFollow(i, j) \leftarrow Follow(i, j)$
 $MayFollow(i, j) \leftarrow Follow(i, k), MayFollow(k, j)$



Follow(1,2)
Follow(2,3)
Follow(3,4)
Follow(2,5)
Follow(5,6)

Datalog input

MayFollow(1,2)
MayFollow(1,3)
MayFollow(1,4)
MayFollow(1,5)
MayFollow(1,6)
MayFollow(2,3)
MayFollow(2,4)
MayFollow(2,5)
MayFollow(2,6)
MayFollow(3,4)
MayFollow(5,6)

Datalog fixpoint

Deriving MayDepOn

```
1: x := Balance
2: Mstore(0x20, x)
3: y := MLoad(0x20)
4: z := x + y
```



```
Follow(1,2)
Follow(2,3)
Follow(3,4)
Assign(x, Balance)
IsConst(0x20)
MStore(2, 0x20, x)
MLoad(3, y, 0x20)
Op(4, z, x)
Op(4, z, y)
```

Derived from
the Balance
instruction

Memory
operations

Capture that
z is derived
from x and y



```
MayDepOn(x, t) ← Assign(x, t)
MayDepOn(x, t) ← Op(⟦, x, x'), MayDepOn(x', t)
MayDepOn(x, t) ← MLoad(l, x, o), isConst(l, o), MemTag(l, o, t)
MayDepOn(x, t) ← MLoad(l, x, o), ¬isConst(l, o), MemTag(l, ⟦, t)
```

```
MemTag(l, o, t) ← MStore(l, o, x), isConst(o), MayDepOn(x, t)
MemTag(l, ⟦, t) ← MStore(l, o, x), ¬isConst(o), MayDepOn(x, t)
MemTag(l, o, t) ← Follows(l, l'), MemTag(l', o, t), ¬MStore(l, o, ⟦)
```

Securify: Under the Hood

```
MemTag(0x20, Balance)  
MemTag(0x40, Const)  
VarTag(z, Const)  
VarTag(k, Gas)  
Assign(s, 0x20)  
Call(s, k)  
⋮
```

Semantic Representation



The screenshot shows the Securify interface with a dark blue background. At the top right, it displays 'TOTAL Issues' with a count of 4. Below this, there are three main categories of issues, each with a count of 1:

- Transaction Reordering**: Includes the issue 'Transactions May Affect Ether Amount' with a description 'The use of concurrency is discouraged in chaincode.' and a 'SimpleBank 10' tag.
- Recursive Calls**: Includes the issue 'Gas-dependent Reentrancy' with a description 'The use of concurrency is discouraged in chaincode.' and a 'SimpleBank 10' tag.
- Insecure Coding Patterns**: Includes the issue 'Unhandled Exception' with a description 'Results of phantom reads should not be used to manipulate the ledger.' and a 'SimpleBank 10' tag.

Each issue entry has an 'info' button and an upward arrow icon.

Patterns DSL

(Labels)	$l ::=$ (labels)
(Vars)	$x ::=$ (variables)
(Tags)	$t ::= l \mid x$
(Instr)	$n ::= \textit{Instr}(l, x, \dots, x)$
(Facts)	$f ::= \textit{MayFollow}(l, l) \mid \textit{MustFollow}(l, l)$ $\mid \textit{MayDepOn}(x, t) \mid \textit{DetBy}(x, t)$
(Patterns)	$p ::= f \mid \forall n: p \mid \exists n: p \mid p \wedge p \mid \neg p$

Detecting the DAO Hack

```
function withdraw() {  
  uint amount = balances[msg.sender];  
  msg.sender.call.value(amount)();  
  balances[msg.sender] = 0;  
}
```

Call instruction
followed by a
write to storage

Formalized as a
trace property

Security property:

No state changes after call instructions

Compliance pattern

$Call(l, _, _, _): \neg \exists SStore(l', _, _). MayFollow(l, l')$

Violation pattern

$Call(l, _, _, _): \exists SStore(l', _, _). MustFollow(l, l')$

Proofs establish a formal logical relation between the property and its patterns

Detecting Unrestricted Writes

```
address owner = ...;  
function initWallet(address _owner) {  
  owner = _owner;  
}
```

Unrestricted
write

Security property:

No storage offset is writable by all users

Formalized as a
hyperproperty

Compliance pattern

$SStore(l, x, _): DetBy(x, Caller)$

Violation pattern

$SStore(l, x, _): \neg MayDepOn(x, Caller)$
 $\wedge \neg MayDepOn(l, Caller)$

The image features a dark blue background with a white grid pattern. A horizontal band of a lighter blue color runs across the middle of the image. The text is centered within this band.

How well does this approach work in practice?

Securify vs. Existing Solutions

State-of-the-art security checkers for Ethereum smart contracts

- Oyente
- Mythril

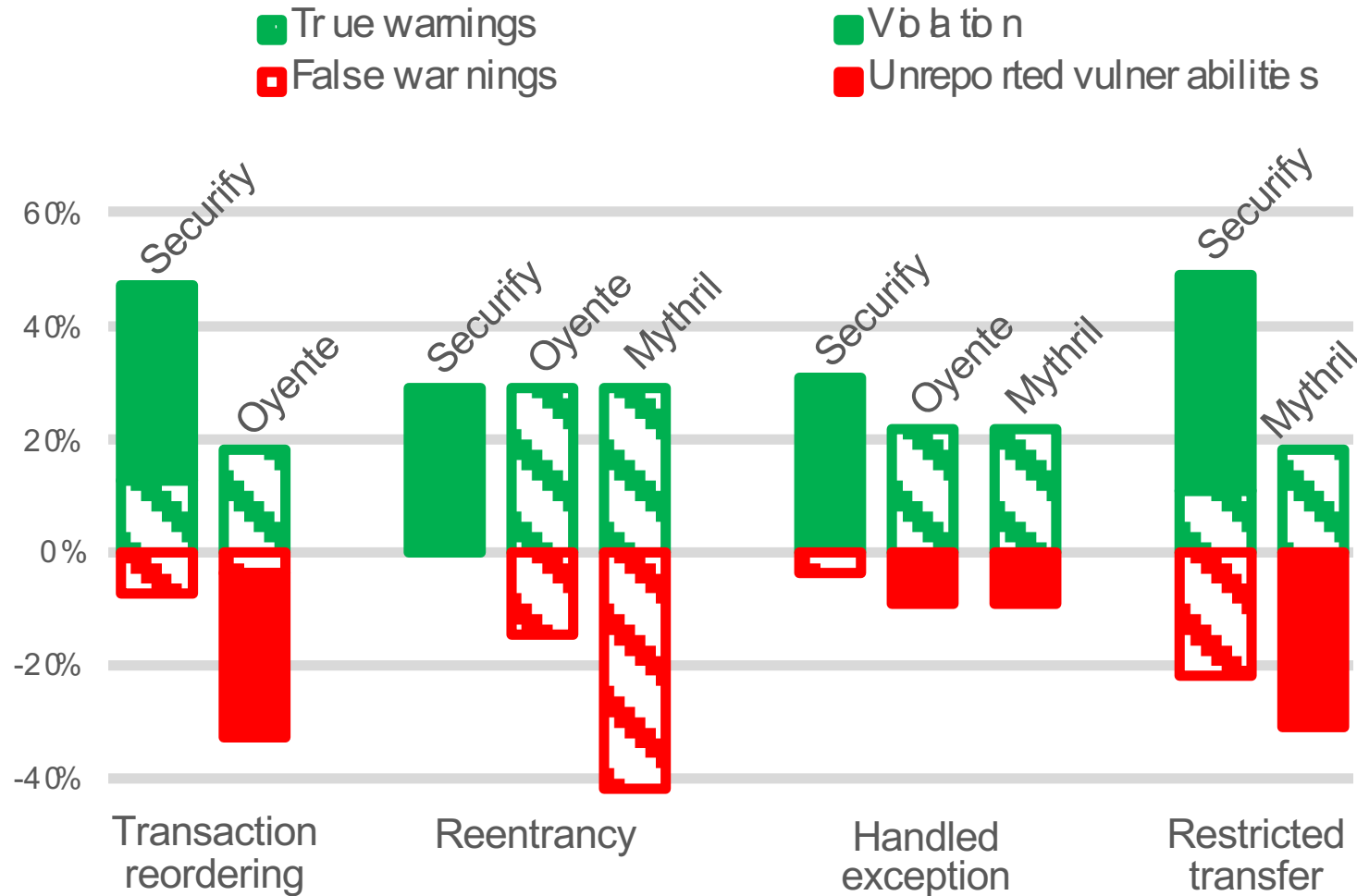
Dataset

- 80 open-source smart contracts

Experiment

- Run contracts using Securify, Oyente, and Mythril
- Manually inspect each reported vulnerability

Securify vs. Oyente vs. Mythril



Research



ICE center@ **ETH**

AI² <http://ai2.ethz.ch>

SECURIFY <http://securify.ch>

DEGUARD <http://apk-deguard.com>

JS NICE <http://isnice.org>

PSI SOLVER <http://psisolver.org>

EVENT RACER <http://eventracer.org>

Start-ups



Securing the blockchain

<https://chainsecurity.com>



jobs@chainsecurity.com



contact@chainsecurity.com



[@chain_security](https://twitter.com/chain_security)